

# IN324: Floyd-Hoare logic



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Audience : IN  
Date :

## 1. Inductive definitions

1. give an inductive definitions of natural numbers
2. give an inductive definitions of binary trees

Prove the following property of binary trees: "the number  $n$  of nodes in a binary tree of height  $h$  is at least  $n = h$  and at most  $n = 2^h - 1$  where  $h$  is the depth of the tree".

## 2. Natural deduction for propositional logic

Prove the following theorems using natural deduction for propositional logic:

(a)  $(a \rightarrow (b \rightarrow c)) \rightarrow ((a \rightarrow b) \rightarrow (a \rightarrow c))$

(b)  $((a \vee b) \rightarrow c) \rightarrow (b \rightarrow c)$

$$(c) \ ((a \vee b) \wedge (a \rightarrow c) \wedge (b \rightarrow c)) \rightarrow c$$

$$(d) \ a \rightarrow \neg\neg a$$



**3. Maths, again and again**

Let  $E$  be a set. Model the following mathematical notions using a first-order language. Define precisely the signature of your language.

(a)  $=$  define the “classical” equality relation on  $E$

(b)  $\leq$  is a preorder on  $E$

(c)  $(E, .)$  is a monoid

**4. Natural deduction for first-order logic**

Prove the following theorems using natural deduction for first-order logic:

(a)  $(\forall x \varphi \wedge \psi) \rightarrow (\forall x \varphi \wedge \forall x \psi)$

(b)  $\exists x \forall y \varphi \rightarrow \forall y \exists x \varphi$

### 5. Floyd-Hoare logic to prove programs

For each of the following program:

- determine the invariant in order to prove the partial correctness of the program
- determine the variant in order to prove the total correctness of the program
- annotate the program using Floyd-Hoare logic

(a) consider the function  $\mathbb{N} \rightarrow \mathbb{N}$  defined by the following program:

```

{N ≥ 0}
K := 0;
F := 1;
while (K ≠ N) do
    K := K + 1;
    F := F * K
od
{F = N!}

```

**Invariant**

**Variant**

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```

{N ≥ 0}

```

```

K := 0;

```

```

F := 1;

```

```

while (K ≠ N) do

```

```

    K := K + 1;

```

```

    F := F * K

```

```

od

```

```

{F = N!}

```

(b) consider the function  $\mathbb{N} \rightarrow \mathbb{N}$  defined by the following program:

|                                       | Invariant | Variant |
|---------------------------------------|-----------|---------|
| $\{N \geq 0\}$                        |           |         |
| $K := N;$                             |           |         |
| $F := 1;$                             |           |         |
| <b>while</b> ( $K \neq 0$ ) <b>do</b> |           |         |
| $F := F * K;$                         |           |         |
| $K := K - 1$                          |           |         |
| <b>od</b>                             |           |         |
| $\{F = N!\}$                          |           |         |

---

$\{N \geq 0\}$

$K := N;$

$F := 1;$

**while** ( $K \neq 0$ ) **do**

$F := F * K;$

$K := K - 1$

**od**

$\{F = N!\}$

(c) consider the function  $\mathbb{N}^2 \rightarrow \mathbb{N}^2$  defined by the following program:

|                                                              | Invariant | Variant |
|--------------------------------------------------------------|-----------|---------|
| $\{X \geq 0 \wedge Y > 0\}$                                  |           |         |
| $Q := 0;$                                                    |           |         |
| $R := X;$                                                    |           |         |
| <b>while</b> $(Y \leq R)$ <b>do</b>                          |           |         |
| $Q := Q + 1;$                                                |           |         |
| $R := R - Y$                                                 |           |         |
| <b>od</b>                                                    |           |         |
| $\{X = Q \times Y + R \wedge 0 \leq Q \wedge 0 \leq R < Y\}$ |           |         |

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$\{X \geq 0 \wedge Y > 0\}$

$Q := 0;$

$R := X;$

**while**  $(Y \leq R)$  **do**

$Q := Q + 1;$

$R := R - Y$

**od**

$\{X = Q \times Y + R \wedge 0 \leq Q \wedge 0 \leq R < Y\}$

(d) consider the function  $\mathbb{N}^2 \rightarrow \mathbb{N}^2$  defined by the following program:

|                                                      | Invariant | Variant |
|------------------------------------------------------|-----------|---------|
| $\{A > 0 \wedge B > 0\}$                             |           |         |
| $X := A;$                                            |           |         |
| $Y := B;$                                            |           |         |
| <b>while</b> $(X \neq Y)$ <b>do</b>                  |           |         |
| <b>if</b> $(X > Y)$ <b>then</b>                      |           |         |
| $X := X - Y$                                         |           |         |
| <b>else</b>                                          |           |         |
| $Y := Y - X$                                         |           |         |
| <b>fi</b>                                            |           |         |
| <b>od</b>                                            |           |         |
| $\{X = Y \wedge X > 0 \wedge X = \text{gcd}(A, B)\}$ |           |         |

where the *gcd* function is defined by:

$$\begin{aligned} \forall A \in \mathbb{N}^* \quad \text{gcd}(A, A) &= A \\ \forall A \in \mathbb{N}^* \forall B \in \mathbb{N}^* \quad \text{gcd}(A, B) &= \text{gcd}(B, A) \\ \forall A \in \mathbb{N}^* \forall B \in \mathbb{N}^* \quad A > B &\rightarrow \text{gcd}(A, B) = \text{gcd}(A - B, B) \end{aligned}$$


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$\{A > 0 \wedge B > 0\}$

$X := A;$

$Y := B;$

**while**  $(X \neq Y)$  **do**

**if**  $(X > Y)$  **then**

$X := X - Y$

**else**

$Y := Y - X$

**fi**



**od**

$$\{X = Y \wedge X > 0 \wedge X = \gcd(A, B)\}$$